

"Tutorial Input"—Standardizing the Computer/Human Interface

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This article describes a new technique for implementing a human/computer interface for computer-based subsystems for the DSN. Known as "tutorial input," this technique provides convenient short input procedures for the experienced operator and a helping hand for the novice. From the programmer's viewpoint, the technique is implemented in a compact, modular, easily modified table-driven structure. The technique has been successfully used through two generations of R&D ranging machines.

I. Motivation

Though much time is spent by programmers in producing efficient, clean, and ego-pleasing code, very little time is spent in developing an efficient, reliable human/computer interface. From an operational standpoint, this interface is the most important and least understood in DSN subsystem programming. With a view toward minimizing operator error and increasing subsystem efficiency, this article will present the technique for human/computer communication successfully used in two generations of R&D ranging systems.

In a typical subsystem program, the operator must provide operating parameters, critical times and perhaps limits to the software. Two interface techniques are generally used. With the first technique the operator may be queried on an input-by-input basis. For instance,¹

¹Computer typeouts are underlined.

ENTER T1: 2

26

ENTER T2: 2

15

ENTER T3: 2

15

A second technique is the preset format wherein the operator must enter numbers in accordance with some specified template. An analogous entry to the one given above could be

*/26/15/15\$

The first technique has the advantage that a format need not be memorized or followed. Given the parameters, an operator just follows directions. The disadvantage is that at the 10 characters per second of the usual tele-

typewriter, the questions sometimes waste an unacceptable amount of time. The second technique is clearly faster but requires a thorough a priori knowledge of the format. This leads to increased training time and a greater chance of operator error from misplaced fields.

Overriding the problems and tradeoffs inherent in these techniques is the multitude of programs involved in a typical DSN operation. An operator may have to communicate with a variety of programs, each designed with a different input philosophy and format. This situation naturally leads to an increased probability of human error.

In a real-time mission environment, an incorrect program entry can be just as disastrous as an equipment malfunction. A specific example occurred to ranging during the Mariner Mars 1971 (MM'71) mission. The software for command and telemetry running on the XDS 920 Telemetry and Command Processors (TCPs) used a dollar sign for a line terminator, whereas the software for the Mu ranging system used a carriage return. On at least two occasions, range data were lost because the operator typed a dollar sign at the end of an input line and walked away thinking his task complete. The ranging software waited the time-out period and then cancelled the necessary input. A simple modification to the software to allow it to recognize both a dollar sign and carriage return as a terminator saved a significant amount of data.

II. Proposal

Usually the subsystem operator is treated as a button-pusher who is taught to run specific software and devices. This view was believed incorrect, and software to support both the Mu-I and Mu-II ranging systems was developed with the direct cooperation and interaction of the operators. The input routine devised in this effort has been shown to be easy to learn as well as easy to operate. It uses a table-driven structure which makes it highly visible to the programmer who must implement it and relatively easy to modify and enlarge. This algorithm, known as "tutorial input," is presented here.

III. Tutorial Input as Viewed by an Operator

There are basically two parts to an input using the "tutorial" technique. First, a command is typed to designate the type of entry; second, the parameter or parameters are entered. For clarity, consider the commands used with the Mu-II system given in Table 1. (Parameter definitions are given in Table 2 for completeness.)

To input commands and data, the operator first notifies the PDP 11/20 that command input is desired by pressing an interrupt button (this could be a breakpoint on the XDS 9-series machines). When the machine responds with a pound sign (#) he types an input line terminated with a carriage return. The first field is a command followed by a slash (/). If the command is acceptable, the software looks to see if anything else is in the input line. If more characters have been typed, they are used. If not, the operator is queried by a specific message for the appropriate parameter of the now-active command. This is true every time a field containing a command or parameter has been used. If more characters are in the input line, they are processed; if not, and more parameters are required, the operator is queried.

The field delimiters used are a slash between a command and parameters and commas between multiple parameters. Note that the whole line is used. This allows commands to be entered contiguously.

The following are equivalent ways to initialize all ranging parameters:

Example 1:

```
#A/ ↵
TOF: ↵
#1200 ↵
SYN FREQ: ↵
# 44.01234 ↵
T1,T2,T3,TC: ↵
# 20,20,20,12 ↵
C1,C2,CN: ↵
# 4,19,3 ↵
MODE: ↵
# - ↵
```

Example 2:

```
#A/1200,44.01234,20,20,20,12,4,19,3,- ↵
```

Example 3:

```
#A/1200,44.01234,20,20,20 ↵
TC: ↵
#12,4,19 ↵
CN: ↵
#3,- ↵
```

Example 1 shows how the operator can be led through the input parameters. In Example 2, an experienced operator has entered all parameters without software prompting. The operator in example 3, after losing his place, has let the computer request the remaining required numbers.

Multiple commands may be entered on a single line. In order to change the mode and number of components, for instance, one could enter

```
#M/124,C/3,10,50 ↵
```

or

```
#M/124,C/ ↵  
C1,C2,CN: ↵  
#3,*,50 ↵
```

The asterisk (*) entered for C2 causes the previous value of C2 to remain unchanged.

A further provision to ease the operator's task is error correction. Input cancellation (control R—R^c) results in immediately exiting the input routine. Character deletion (C^c) causes the program to type ← and results in the deletion of the last character entered into the string. Character deletion may be used more than once, e.g., to delete the last three characters so that they can be re-typed. Line deletion (E^c) deletes the line, upspaces, types →, and allows the whole line to be re-typed. Error correction can also be accomplished by use of the command to change a particular parameter.

IV. Tutorial Input as Viewed by a Programmer

Although written in machine language for a PDP-11/20, via the SAPDP Xerox Sigma 5 cross assembler (Ref. 1) the routine is amenable to coding in another machine language or in a higher level language such as BASIC or FORTRAN. In this discussion the interaction of the input routine with other real-time processors will not be covered. The Mu-II software required interfaces to a teletype output routine and to a real-time scheduler. These topics will be documented in a forthcoming article on the Mu-II system software.

Though perhaps not meeting the letter of structured programming, which is difficult, if not impossible, in a real-time environment, tutorial input does realize all advantages normally claimed for structured programming. Through the use of a table-driven technique, the input

sequencing and programming is straightforward, easily modified, and simply documented.

The technique involves two tables. The first table is the command list. Each entry in the list takes three words in the Mu-II realization. There is one entry for each command. The first word has the two characters of the command, e.g., "A/" or "TF," the second a pointer to a word in the second table, and the third word a number equal to the number of parameters to be entered with the particular command.

The second table is the parameter list. Each entry in this list takes four words. There is one entry for each parameter. The first word of the entry contains the number of characters in the message associated with each parameter, while the second contains the message location. (The messages are the queries shown earlier, e.g., TOF: ↵.) The third word gives the location of the decoding subroutine (integer, floating point, etc.) to convert the parameter's ASCII character string to binary. The fourth and final word contains the location of the destination for the binary number. Figure 1 contains a listing of the two tables in the Mu-II software.

The algorithm which interprets these tables is described structurally in Fig. 2 and in a detailed flow chart in Fig. 3.

As an example of the algorithm, let us consider that a command "T/" has been input and follow the algorithm operation. Scanning the command list (Fig. 1), a match is found (line 334) giving the parameter table location as SCT1 and the number of parameters to be input as four. The program looks to see if there are any more characters in the input string. If not, the message TMSG, containing 14 characters, is typed as shown by the SCT1 entry in the parameter list. If input already exists, the message is skipped. In either case, the input parameter is processed by INGR which is the integer decoding subroutine and the result stored in IT1.

Thus far, one parameter has been processed and three more are left. The program proceeds directly to the SCT2 entry in the parameter table and repeats the input process. This continues until all parameters are entered.

Once again the algorithm checks to see if there are more characters in the input stack. If there are none, the routine exits. If there are more, the first field is used to search the command stack as the program assumes that another command is in the input string. The process goes on until an error occurs or until the input line is exhausted.

The input logic flow is readily observable by following the driving tables. The tables are also a simple documentation of the input sequences. The ease with which commands and entries can be modified or deleted is shown by a program change that occurred during the Mariner Venus/Mercury 1973 (MVM'73) mission.

As originally written, an operator could modify any parameter or set of parameters individually except for time of flight (TOF). The only way to change this parameter was through the "enter all parameters" A/ command. As the mission proceeded, the TOF changed by several seconds each day. It was obvious that going through the entire initialization sequence to change TOF was a waste of time. The decision was made to add the TF/ command so that TOF could be changed individually.

Because of the table-driven structure, only three cells were actually needed. These were, from Fig. 1, lines 346-348:

"TF"	Command
SCTOF	Parameter Table Pointer
1	Number of parameters to be entered.

This simple modification changed a tedious operation into a trivial one.

V. Summation

Data will continue to be lost due to incompatibilities, indeed contradictions, between input formats in various DSN software systems. This coupled with the trend toward smaller station operational crews, automation, and the continued proliferation of minicomputers in the DSN makes standardization increasingly important. The algorithm presented herein is a candidate to aid in that task.

A final comment, tutorial input has been used in ranging software throughout the MVM'73 mission. It has shown itself to be an easy-to-use as well as an easy-to-learn input technique. The extra time taken to interact with experienced ranging operators during its development has been paid back many fold through simplified training and the low probability of human error. This programming effort clearly shows the value of bringing the system operators into the software design process at a very early stage. Many times the operators were able to point out features which were not helpful and could be discarded as well as request features which would simplify their task. The success of the Mu-II system programming is largely due to this cooperation between the programmer and the system operators.

Reference

1. Erickson, D. E., "The SAPDP Program Set for Sigma 5 Assembly," in *The Deep Space Network Progress Report*, Technical Report 32-1526, Vol. VII, pp. 91-96, Jet Propulsion Laboratory, Pasadena, Calif., Feb. 15, 1972.

Table 1. Mu-II commands

Command	Parameter to be entered
A/	All operational parameters: TOF, SYN FREQ, T1,T2,T3,TC,C1,C2,CN,MODE
TF/	TOF
C/	C1,C2,CN
T/	T1,T2,T3,TC
S/	SYN FREQ
M/	MODE
Z/	Requested TØ time
Y/	Typewriter printout ON/OFF: i.e., Y/ON Y/OFF

Table 2. Parameter definitions

Abbreviation	Meaning
TOF	Round-trip light time
SYN FREQ	Exciter synthesizer frequency
T1	First component integration time
T2	Lower-frequency components integration time
T3	Post-acquisition DRVID integration time
TC	10-MHz calibration integration time
C1	Highest-frequency component
C2	Lowest-frequency component
CN	Number of post-acquisition DRVID points before automatic reacquisition
MODE	Select configuration (i.e., Block III or IV receiver phasing, reac- quisition with or without coder sync, etc.) “-” results in standard configuration
TØ	Coder synchronization time

FOO	14151	APR 05, '74					SUPER MU II
326*						PAGE	
327*						*COMMAND LIST	
328*	01	007B2	2	C1	A	CONSCAN BYTE	01301,01257 A/
	01	007B2	3	AF	A		
329*	01	007B3	1	EFA	N	WORD	SCT0F
330*	01	007B3	2	000A	A	WORD	10
331*	01	007B4	C3		A	BYTE	01303,01257 C/
	01	007B4	1	AF	A		
332*	01	007B4	2	1F2A	N	WORD	SCC1
333*	01	007B5	0003		A	WORD	3
334*	01	007B5	2	D4	A	BYTE	01324,01257 T/
	01	007B5	3	AF	A		
335*	01	007B6	1	F0A	N	WORD	SCT1
336*	01	007B6	2	0004	A	WORD	4
337*	01	007B7	D3		A	BYTE	01323,01257 S/
	01	007B7	1	AF	A		
338*	01	007B7	2	1F02	N	WORD	SCSYNF
339*	01	007B8	0001		A	WORD	1
340*	01	007B8	2	DA	A	BYTE	01332,01257 Z/
	01	007B8	3	AF	A		
341*	01	007B9	1	F4A	N	WORD	SCT0
342*	01	007B9	2	0001	A	WORD	1
343*	01	007BA	CD		A	BYTE	01315,01257 M/
	01	007BA	1	AF	A		
344*	01	007BA	2	1F42	N	WORD	SCMDE
345*	01	007BB	0001		A	WORD	1
346*	01	007BB	2	D4	A	BYTE	01324,01306 TF
	01	007BB	3	C6	A		
347*	01	007BC	1	EFA	N	WORD	SCT0F
348*	01	007BC	2	0001	A	WORD	1
349*	01	007BD	D3		A	CONEND BYTE	01331,01257 Y/
	01	007BD	1	AF	A		
350*	01	007BD	2	1F52	N	WORD	TYCMD
351*	01	007BE	0001		A	WORD	1
352*						PAGE	
353*						*PARAMETER LIST	
354*	01	007BE	2	0006	A	SCT0F WORD	6
355*	01	007BF	1	F5A	N	WORD	T0FMSG
356*	01	007BF	2	1BFO	N	WORD	INGR
357*	01	007C0	2	2872	N	WORD	IT0F
358*	01	007C0	2	000E	A	SCSYNF WORD	14
359*	01	007C1	1	F60	N	WORD	SYNFMMSG
360*	01	007C1	2	1C3E	N	WORD	INDPFP
361*	01	007C2	2	2874	N	WORD	ISYNF
362*	01	007C2	2	000E	A	SCT1 WORD	14
363*	01	007C3	1	F6E	N	WORD	TMSG
364*	01	007C3	2	1BFO	N	WORD	INGR
365*	01	007C4	2	2864	N	WORD	IT1
366*	01	007C4	2	0006	A	SCT2 WORD	6
367*	01	007C5	1	F80	N	WORD	T2MSG
368*	01	007C5	2	1BFO	N	WORD	INGR
369*	01	007C6	2	2866	N	WORD	IT2
370*	01	007C6	2	0006	A	SCT3 WORD	6
371*	01	007C7	1	F86	N	WORD	T3MSG
372*	01	007C7	2	1BFO	N	WORD	INGR
373*	01	007C8	2	2868	N	WORD	IT3
374*	01	007C8	2	0006	A	SCTC WORD	6
375*	01	007C9	1	F8C	N	WORD	TCMSG
376*	01	007C9	2	1BFO	N	WORD	INGR
377*	01	007CA	2	286A	N	WORD	ITC
378*	01	007CA	2	000E	A	SCC1 WORD	14
379*	01	007CB	1	F7C	N	WORD	CMSG
380*	01	007CB	2	1BFO	N	WORD	INGR
381*	01	007CC	2	286C	N	WORD	IC1
382*	01	007CC	2	0006	A	SCC2 WORD	6
383*	01	007CD	1	FC2	N	WORD	C2MSG
384*	01	007CD	2	1BFO	N	WORD	INGR
385*	01	007CE	2	286E	N	WORD	IC2
386*	01	007CE	2	0006	A	SCCN WORD	6
387*	01	007CF	1	FC8	N	WORD	CNMSG
388*	01	007CF	2	1BFO	N	WORD	INGR
389*	01	007D0	2	286E	N	WORD	ICN
390*	01	007D0	2	000A	A	SCMDE WORD	10
391*	01	007D1	1	F8A	N	WORD	MMSG
392*	01	007D1	2	1CC0	N	WORD	INMDE
393*	01	007D2	2	2870	N	WORD	IMDE
394*	01	007D2	2	000E	A	SCT0 WORD	14
395*	01	007D3	1	F94	N	WORD	T0MSG
396*	01	007D3	2	1D60	N	WORD	INT0
397*	01	007D4	2	288C	N	WORD	ITO
398*	01	007D4	2	000E	A	TYCMD WORD	14
399*	01	007D5	1	FA2	N	WORD	TYMSG
400*	01	007D5	2	1C92	N	WORD	CNTRL
401*	01	007D6	2	287C	N	WORD	ITYSWT

Fig. 1. Mu-II software listing

```

Command, IF command request DO.
    Save current parameters.
    Input line.
    DO UNTIL input stack empty.
        Separate next input field.
        Search command table for match.
        IF NOT match.
            Notify operator.
            EXIT command.
        ELSE.
            Location ← command table (match +1).
            Count ← command table (match +2).
        Parameter, DO UNTIL count = 0.
            If input stack NOT EMPTY.
                Separate next input field.
            ELSE
                Print parameter table (location) characters from
                    string addressed by parameter table (location +1).
                Input line.
                Location ← location +2.
                DO SUBROUTINE addressed by parameter table (location).
                Location ← location +1.
                Store result in parameter addressed by parameter table (location)
                Location ← location +1.
                Count ← count -1
            END
        END
    END
END

```

Fig. 2. Structured representation

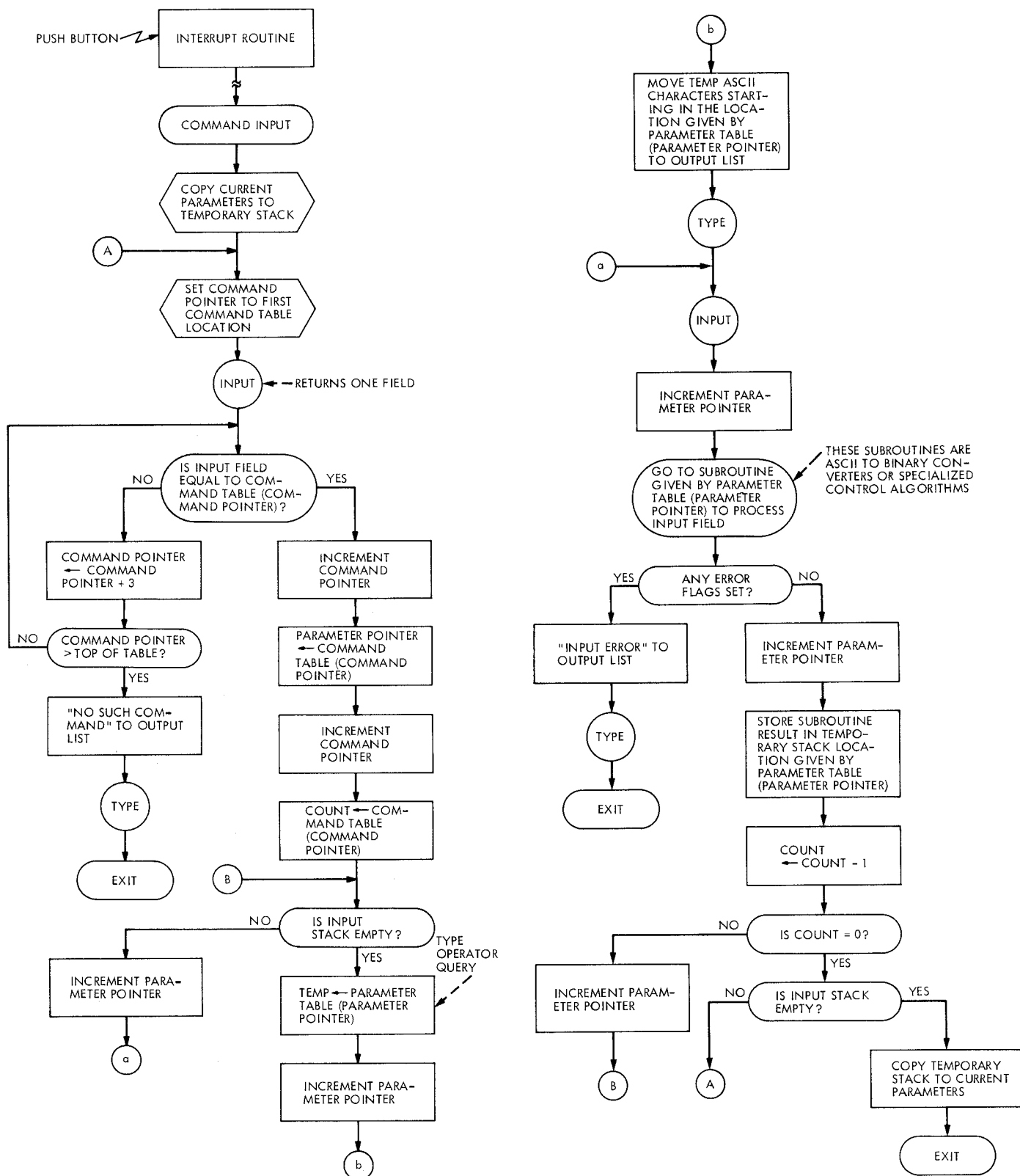
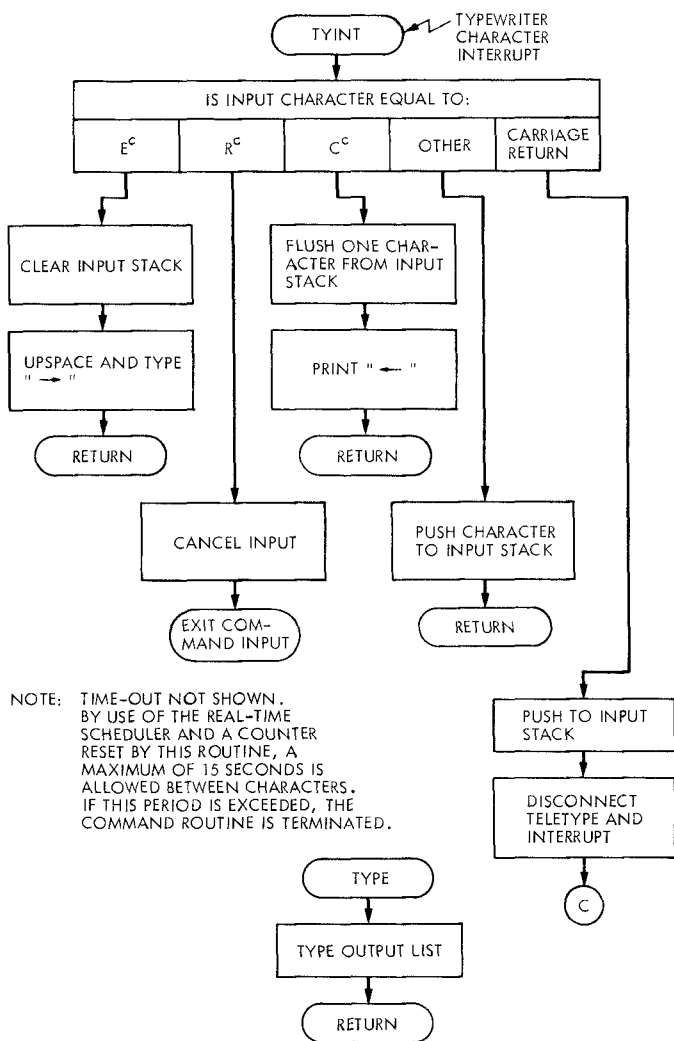
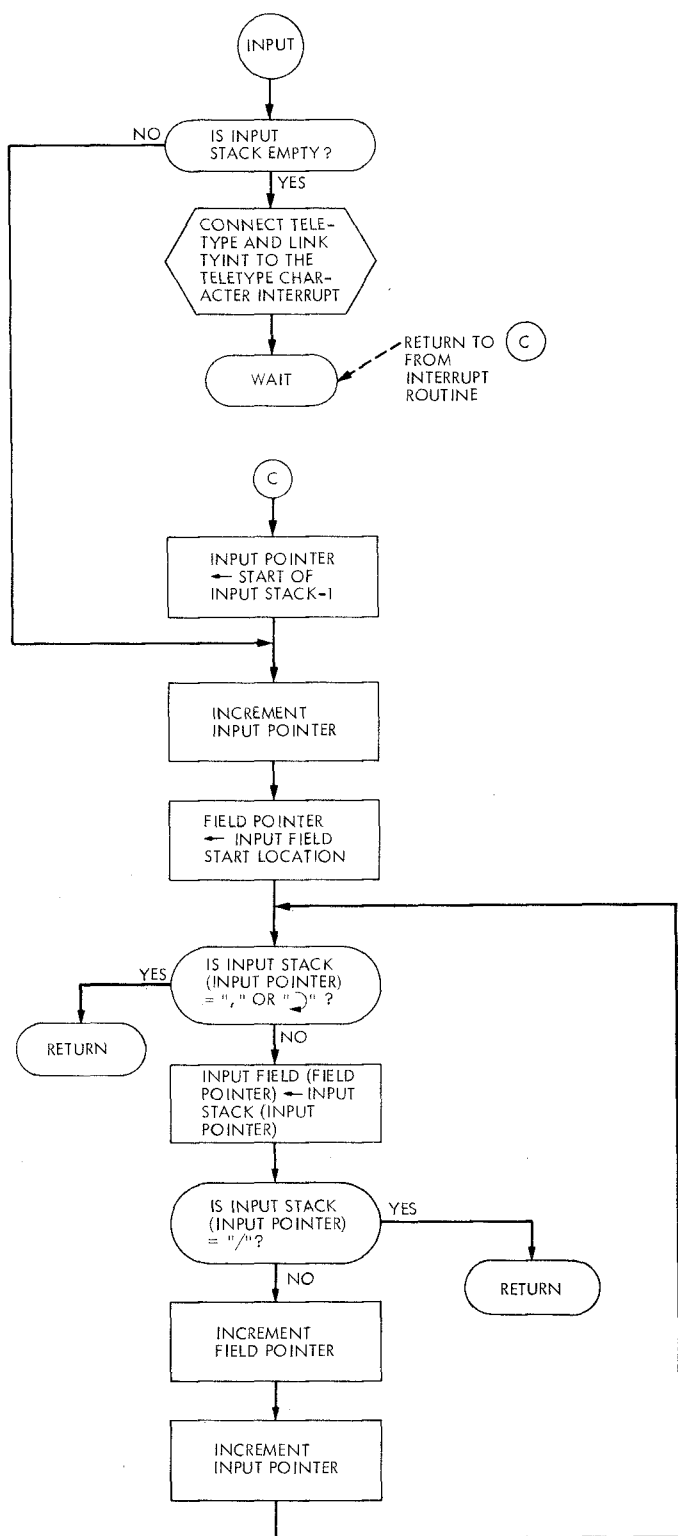


Fig. 3. Tutorial input schematic flowchart (Mu-II realization)



TABLES

COMMAND TABLE		PARAMETER TABLE	
"COMMAND" -ASCII		NO. OF MESSAGE CHARACTERS	
POINTER		MESSAGE LOCATION	
NUMBER OF PARAMETERS		DECODING SUBROUTINE	
		TARGET LOCATION	

Fig. 3 (contd)